Can the quantum switch be deterministically simulated?

Jessica Bavaresco, Satoshi Yoshida, Tatsuki Odake, Hlér Kristjánsson, Philip Taranto, Mio Murao, Marco Túlio Quintino

Sorbonne Université, CNRS, LIP6

2nd FoQaCiA Workshop 03/06/2025







Talk based on:

1. arXiv:2409.18420 [pdf, other] quant-ph

Exponential separation in quantum query complexity of the quantum switch with respect to simulations with standard quantum circuits

Authors: Hlér Kristjánsson, Tatsuki Odake, Satoshi Yoshida, Philip Taranto, Jessica Bavaresco, Marco Túlio Quintino, Mio Murao

Abstract: Quantum theory is consistent with a computational model permitting black-box operations to be applied in an indefinite causal order, going beyond the standard circuit model of computation. The quantum switch — the simplest such example — has been shown to provide numerous information-processing advantages. Here, we prove that the action of the quantum switch on two n-qubit quantum channels can... ¬ More

Submitted 1 October, 2024; v1 submitted 26 September, 2024; originally announced September 2024.

Comments: 23 pages, 3 figures

2. arXiv:2409.18202 [pdf, other] quant-ph

Can the quantum switch be deterministically simulated?

Authors: Jessica Bavaresco, Satoshi Yoshida, Tatsuki Odake, Hlér Kristjánsson, Philip Taranto, Mio Murao, Marco Túlio Quintino

Abstract: Higher-order transformations that act on a certain number of input quantum channels in an indefinite causal order - such as the quantum switch - cannot be described by standard quantum circuits that use the same number of calls of the input quantum channels. However, the question remains whether they can be simulated, i.e., whether their action on their input channels can be deterministically repr... ∇ More

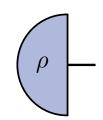
Submitted 26 September, 2024; originally announced September 2024.

Comments: 16 + 14 pages, 4 + 5 figures

Formalism: Higher-order operations

quantum data:

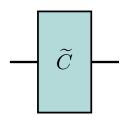
quantum states

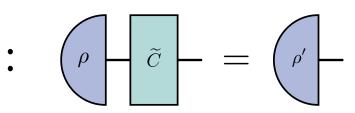


$$\rho \in \mathcal{L}(\mathcal{H}_{\mathrm{in}})$$

quantum functions:

quantum operations (quantum channels)

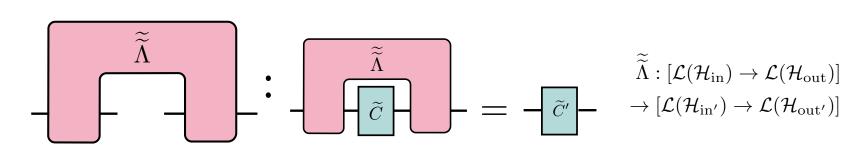




$$\widetilde{C}: \mathcal{L}(\mathcal{H}_{\mathrm{in}}) \to \mathcal{L}(\mathcal{H}_{\mathrm{out}})$$

higher-order quantum operations:

quantum processes



"functions of functions"

Key features: Higher-order operations

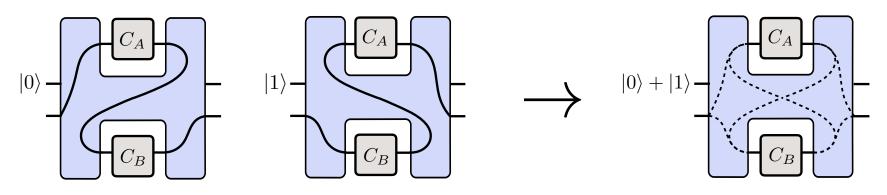
2) Quantum higher-order operations go beyond the quantum circuit model

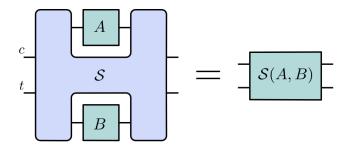
CLASSICAL	Quantum	Higher-order Quantum	
data (bits)	quantum data (qbits)	quantum data (qbits)	
gates	quantum gates	quantum gates	
circuits	circuits	quantum circuit structure	

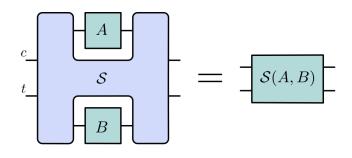
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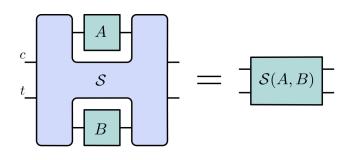
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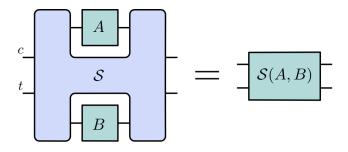




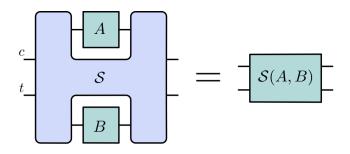
If
$$A(
ho)=U_A
ho U_A^\dagger$$
 and $B(
ho)=U_B
ho U_B^\dagger$



If
$$A(\rho)=U_A\rho U_A^\dagger$$
 and $B(\rho)=U_B\rho U_B^\dagger$
 $\mathcal{S}:(U_A,U_B)\mapsto |0\rangle\!\langle 0|\otimes U_BU_A+|1\rangle\!\langle 1|\otimes U_AU_B$



What can we do with that?



What can we do with that?

The commuting, anti-commuting game:
Perfect discrimination of no-signalling channels via quantum superposition of causal structures

G. Chiribella, PRA 2012

Witnessing causal nonseparability

M. Araújo, C. Branciard, F. Costa, A. Feix, C. Giarmatzi, Č. Brukner, NJP 2015

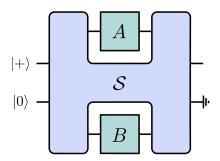
 (U_A, U_B) is a pair of unitary that

$$U_A U_B = U_B U_A$$
 or $U_A U_B = -U_B U_A$

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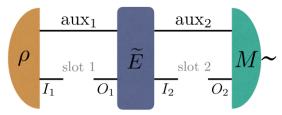
The quantum switch is useful:



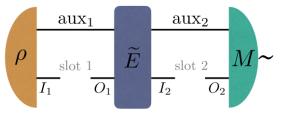
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- Standard ordered strategy:



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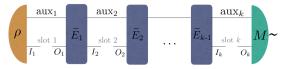
▶ We can find finite sets of unitaries such that $p_{\text{ordered}} \leq 0.87$.

► Great, the quantum switch is useful!

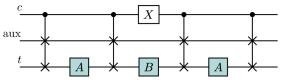
- ► Great, the quantum switch is useful!
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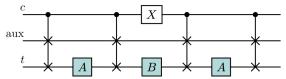
- ► Great, the quantum switch is useful!
- ▶ But . . .
- How big is this advantage? What if we do not have the quantum switch, but we have access to more queries?
- ▶ With a single extra query, sequential strategies can decide if (U_A, U_B) commutes or anti-commutes



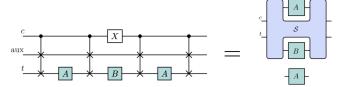
Quantum computations without definite causal structure
 G. Chiribella, G. M. D'Ariano, P. Perinotti, B. Valiron, PRA (2013)



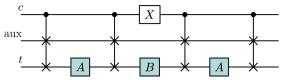
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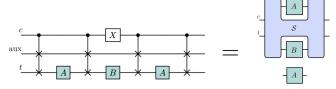
► If A and B are unitary:



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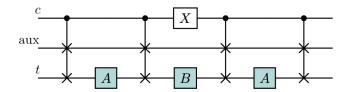


The switch is essentially useless for query complexity tasks...

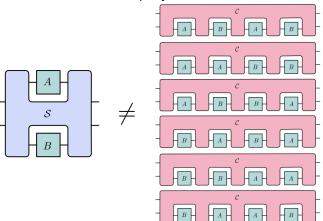
► Wait...

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- What if the operations are not unitary?

- ► Wait...
- ▶ What if the operations are not unitary?
- ► E.g., $A(\rho) = B(\rho) = \operatorname{tr}(\rho) \frac{\mathbb{I}}{d}$

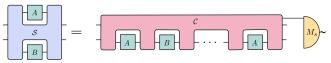


Thm1: There is no quantum circuit that simulates the quantum switch when one extra query of each channel is available.

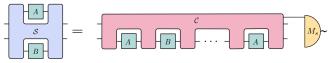


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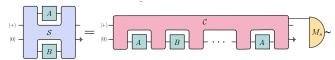
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- Probabilistic simulation:



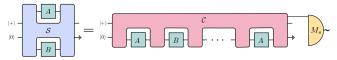
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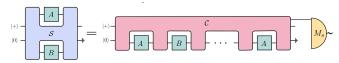
Restricted probabilistic simulation:



How about the probabilities?

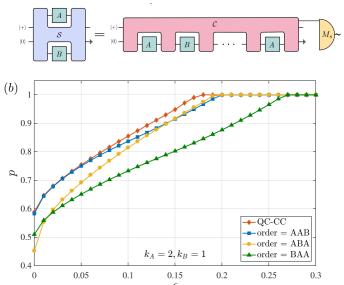


► How about the probabilities?

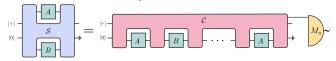


(k_A,k_B)	order	probability
(1, 1)	AB	$p < \frac{4001}{10000}$
	AAB	$p < rac{5715}{10000}$
(2,1)	ABA	$p < \frac{4919}{10000}$
	BAA	$p < \frac{5001}{10000}$

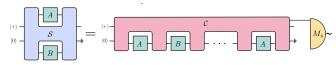
▶ The result is also robust, $F(S, S_{sim}) = 1 - \epsilon$



▶ How about the probabilities when k = 4?

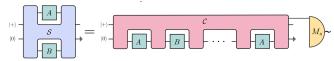


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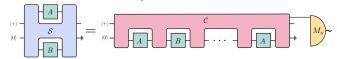
	AABB	$p < \frac{8307}{10000}$
(2, 2)	ABAB	$p < \frac{8484}{10000}$
	ABBA	$p < \frac{8695}{10000}$
	AAAB	$p < rac{8373}{10000}$
(3, 1)	AABA	$p < \frac{6909}{10000}$
	ABAA	$p < rac{7597}{10000}$
	BAAA	$p < rac{6845}{10000}$

► How about identical channels?



k	order	probability
2	AA	$p < \frac{4001}{10000}$
3	AAA	$p < \frac{6534}{10000}$
4	AAAA	$p = 1 \; (*)$

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If the target is not discarded, p < 0.942 for the AAAA order

► How about unitary channels?

(k_A,k_B)	order (unitary only)	probability
(1, 1)	AB	$p \approx 0.400$
	AAB	$p \approx 0.596$
(2, 1)	ABA	p=1
	BAA	$p\approx 0.607$
	AABB	$p = 1 \ (*)$
(2,2)	ABAB	p = 1
	ABBA	p=1
(3, 1)	AAAB	$p \approx 0.708$
	AABA	p=1
	ABAA	p=1
	BAAA	$p = 1 \ (*)$

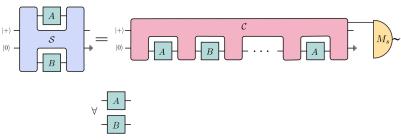
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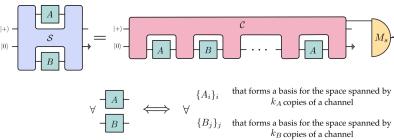
If the target is not discarded, p < 0.822 for the AABB order and p < 0.667 for the BAAA order

► How these results were obtained?

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- ► Optimise over all inputs:



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► SDP (using splitting conic solver)

$$\begin{split} & \textbf{given} & \ \{J_i^A\}_i, \{J_j^B\}_j, k_A, k_B \\ & \textbf{max} \ p \\ & \textbf{s.t.} & \ C_s * [(J_i^A)^{\otimes k_A} \otimes (J_j^B)^{\otimes k_B}] = p \, S * (J_i^A \otimes J_j^B) \ \ \forall i,j \\ & \ C_s \geq 0, \ \ C - C_s \geq 0, \\ & \ \mathbb{P}(C) = C, \ \ \text{tr}(C) = d_{c_I} d_{t_I} d_{AO}^{k_A} d_{BO}^{k_B}, \end{split}$$

$$\begin{split} & \textbf{given} \quad \{J_i^A\}_{i*} \{J_j^B\}_{j}, k_A, k_B \\ & \textbf{min} \quad \frac{1}{d_{A_i^A}^{A_i} d_{B_i^B}^{B_i} d_{c_O} d_{t_O}} \operatorname{tr}(\Gamma) \\ & \textbf{s.t.} \quad \sum_{i,j} \operatorname{tr} \left[R_{ij} \left(S* \left(J_i^A \otimes J_j^B\right)\right)\right] = 1 \\ & \Gamma - \sum_{i,j} R_{ij} \otimes \left[\left(J_i^A\right)^{\otimes k_A} \otimes \left(J_j^B\right)^{\otimes k_B}\right]^T \geq 0 \\ & \Gamma \geq 0, \quad \overline{\mathbb{P}}(\Gamma) = \Gamma. \end{split}$$

any feasible point that yields some $p<1 \label{eq:point}$ constitutes a valid upper bound

▶ But...is that a mathematical proof?

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- ▶ No! But, we can extract a proof out of it!

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Algorithm:

- 1. Construct symbolic non-floating point operators $\Gamma^{\rm sym}$ and $R^{\rm sym}_{ij}$ by truncating them and obtaining a symbolic operator with only rational numbers.
- 2. Force the operators Γ^{sym} and R^{sym}_{ij} to be self-adjoint by making use of the expression $(M+M^\dagger)/2$, which is self-adjoint for any M.
- $\begin{array}{l} 3. \; \text{Evaluate} \; t^{\text{sym}} := \sum_{i,j} \operatorname{tr} \left[R_{ij}^{\text{sym}} \left(S * (J_i^A \otimes J_j^B) \right) \right], \\ \text{where} \; S, J_i^A, \; \text{and} \; J_j^B \; \; \text{are also symbolic} \\ \text{operators.} \; \; \text{Define} \; R_{ij}^{\text{ok}} := \; R_{ij}^{\text{sym}} / t^{\text{sym}} \; \text{for all} \\ i.j. \end{array}$

- 4. Project Γ^{sym} onto the appropriate subspace to obtain $\overline{\mathbb{P}}(\Gamma^{\text{sym}})$.
- 5. Find $\eta \in \mathbb{R}$ such that $\Gamma^{\mathrm{ok}} := \overline{\mathbb{P}}(\Gamma^{\mathrm{sym}}) + \eta \mathbb{1} \geq 0$ and $\Gamma^{\mathrm{ok}} \sum_{i,j} R^{\mathrm{ok}}_{ij} \otimes (J^{A^{\otimes k_A}}_i \otimes J^{B^{\otimes k_B}}_j)^T \geq 0$
- 6. Output the quantity ${\rm tr}(\Gamma^{\rm ok})/d_{c_I}d_{t_I}d_{A_O}^{k_A}d_{B_O}^k$, which is a rigorous upper bound of the primal problem.

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- ▶ Thm2: Let $d = 2^n$ be the dimension of the target state. If $k_A = 1$ and $k_B < 2^n$, there is no quantum circuit simulation of the quantum switch.

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 We then analyse the case of Pauli operations using a differentiation technique (from Analytical lower bound on query complexity for transformations of unknown unitary operations, T. Odake, S. Yoshida, M. Murao).

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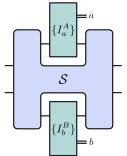
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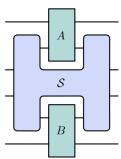
In particular, its eigendecomposition cannot be compatible with QCQC processes.

► How to go beyond?

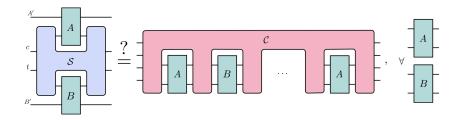
- ► How to go beyond?
- ▶ The quantum switch is not restricted to single-partite channels.

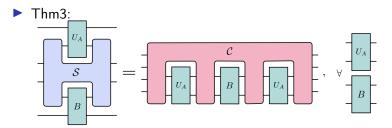
- ► How to go beyond?
- ▶ The quantum switch is not restricted to single-partite channels.
- How about instruments, bipartite channels?



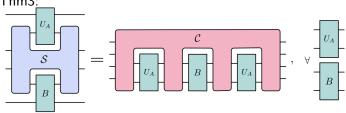


General simulation:

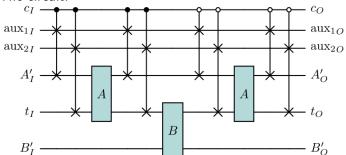




► Thm3:



► The circuit:



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- Open problem: Prove/disprove the conjecture, analyse the scaling (how the probability scales, how the error scales)
- ▶ Dream: "Simple task" that is easy if we have the quantum switch, but very hard if we only have standard circuits

Thank you

