Parallel, sequential, and non-causal strategies for transforming unitary operations and discriminating quantum channel via a higher-order approach

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$$|\psi_{\mathsf{in}}
angle \mapsto |\psi_{\mathsf{out}}
angle$$

$$|\psi_{\mathsf{in}}
angle \mapsto |\psi_{\mathsf{out}}
angle = U|\psi_{\mathsf{in}}
angle$$

$$U^{\dagger}U=I$$

$$ho_{\mathsf{in}} \mapsto 
ho_{\mathsf{out}}$$

$$ho_{\mathrm{in}}\mapsto
ho_{\mathrm{out}}=\widetilde{\Lambda}(
ho_{\mathrm{in}})$$

$$\widetilde{\Lambda}$$
 is CPTP

## Quantum operation transformations

Can we transform quantum operations??

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$$U_{\mathsf{in}}\mapsto U_{\mathsf{out}}$$
  $\widetilde{\Lambda_{\mathsf{in}}}\mapsto \widetilde{\Lambda_{\mathsf{out}}}$ 

"Quantum" unitary inversion

$$U_d \mapsto U_d^{-1}$$

## The universal/unknown paradigm

$$\sigma_{Z} = \begin{bmatrix} 1 & 0 \\ 0 & -1 \end{bmatrix} = \sigma_{Z}^{-1}$$

Ideally...

Something like this:

$$-\sigma_Y - U_2 - \sigma_Y - = -U_2^* -$$

Phys. Rev. Research (2019) J. Miyazaki, A. Soeda, and M. Murao

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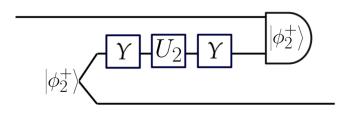
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- Optimal average fidelity:  $F_{max} = \frac{2}{d^2}$ G. Chiribella and D. Ebler, New Journal of Physics (2016)
- $ightharpoonup F_{max} < 1 \implies \mathsf{Impossible...}$

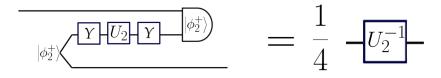
Probabilistic heralded?

Probabilistic heralded? For qubits, Possible!

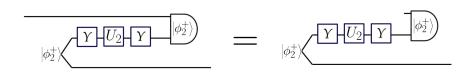
## Explicit construction



# Explicit construction



## Delayed input state protocols

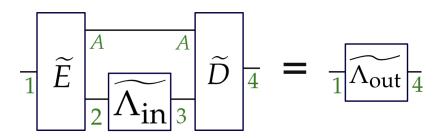


► Is it optimal?

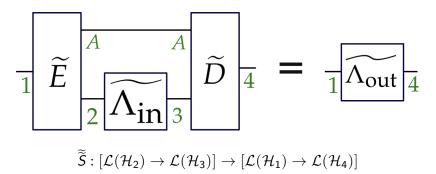
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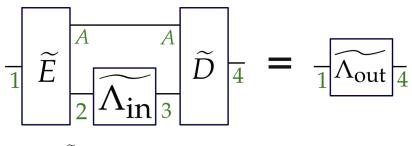
- ► Is it optimal?
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- How can we increase the success probability?
- ► Higher-order operations and supermaps!



The most general quantum superchannel?

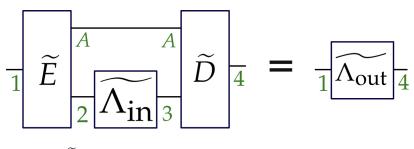


 $ightharpoonup \widetilde{\widetilde{S}}$  is a linear supermap



$$\widetilde{\widetilde{S}}: [\mathcal{L}(\mathcal{H}_2) \to \mathcal{L}(\mathcal{H}_3)] \to [\mathcal{L}(\mathcal{H}_1) \to \mathcal{L}(\mathcal{H}_4)]$$

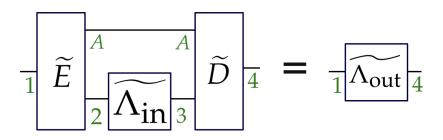
- $\triangleright \tilde{S}$  is a linear supermap
- $\tilde{S}$  maps valid channels into valid channels (TP preserving, CP preserving)



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- $\triangleright$   $\tilde{S}$  is a linear supermap
- $\widetilde{S}$  maps valid channels into valid channels (TP preserving, CP preserving)
- $\triangleright$   $\tilde{S}$  may be applied into part of channel (completely CP preserving)

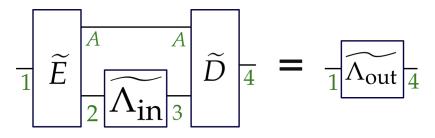




$$\widetilde{\widetilde{S}}(\widetilde{\Lambda_{\mathsf{in}}}) = \mathsf{tr}_{\mathcal{A}}\left(\widetilde{D} \circ \left(\widetilde{\Lambda_{\mathsf{in}}} \otimes \widetilde{\mathit{I}_{\mathcal{A}}}\right) \circ \widetilde{\mathit{E}}\right)$$

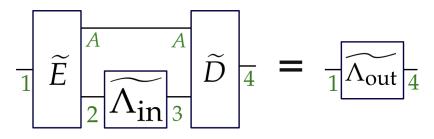
- G. Chiribella, G. M. D'Ariano, and P. Perinotti EPL (2008)
- K. Życzkowski J. Phys. A 41, 355302-23 (2008)
- G. Gutoski and J. Watrous Proceedings of STOC (2007)





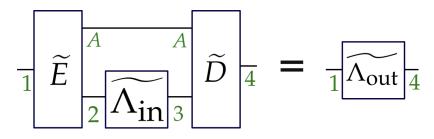
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How to represent such mathematical objects?



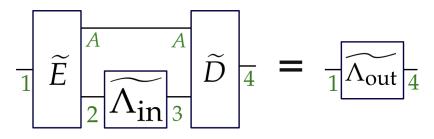
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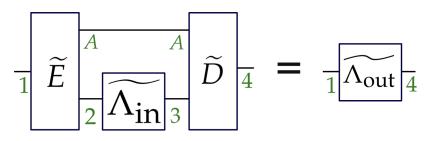
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- ▶ Maps  $\widetilde{\Lambda}$  :  $\mathcal{L}(\mathcal{H}_2) \to \mathcal{L}(\mathcal{H}_3)$  become matrices  $\Lambda \in \mathcal{L}(\mathcal{H}_2 \otimes \mathcal{H}_3)$ .



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- ▶ Supermaps  $\widetilde{\widetilde{S}}$  become maps  $\widetilde{S}: \mathcal{L}(\mathcal{H}_2 \otimes \mathcal{H}_3) \to \mathcal{L}(\mathcal{H}_1 \otimes \mathcal{H}_4)$



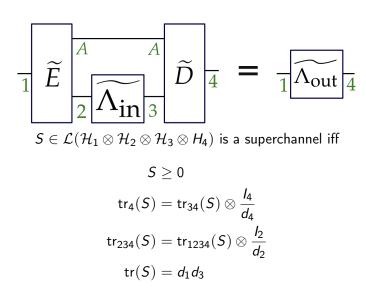


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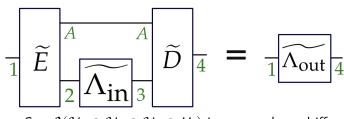
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- ▶ Maps  $\widetilde{S}$  become matrices  $S \in \mathcal{L}(\mathcal{H}_1 \otimes \mathcal{H}_2 \otimes \mathcal{H}_3 \otimes \mathcal{H}_4)$



# Superchannels



# Superchannels



 $S \in \mathcal{L}(\mathcal{H}_1 \otimes \mathcal{H}_2 \otimes \mathcal{H}_3 \otimes \mathcal{H}_4)$  is a superchannel iff

$$S \geq 0$$

$$\mathsf{tr}_4(S) = \mathsf{tr}_{34}(S) \otimes \frac{I_4}{d_4}$$

$$\mathsf{tr}_{234}(S) = \mathsf{tr}_{1234}(S) \otimes \frac{I_2}{d_2}$$

$$\mathsf{tr}(S) = d_1 d_3$$

Affine and positive semidefinite constraints  $\implies$  SDP!!



► Is it optimal?

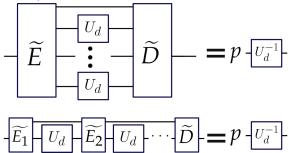
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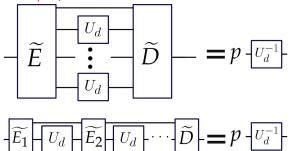
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- ► More calls/copies

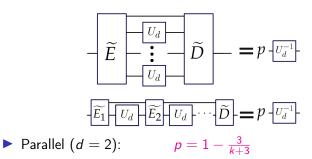


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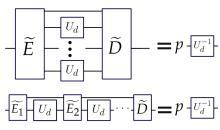


(Quantum combs, channel with memory, quantum strategy, quantum superchannels with multiple inputs)



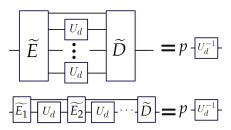






- ▶ Parallel (d = 2):  $p = 1 \frac{3}{k+3}$
- Parallel (k < d 1): p = 0

$$o = 0$$



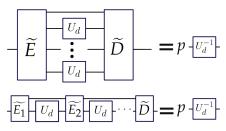
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$$p=1-\tfrac{3}{k+3}$$

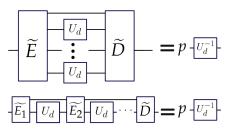
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$$p = 0$$

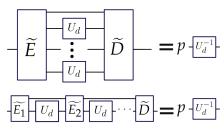
▶ Parallel 
$$(k \ge d-1)$$
:  
 $1 - \frac{1}{k} \sim 1 - \frac{d^2 - 1}{\left\lfloor \frac{k}{d-1} \right\rfloor + d^2 - 1} \le p \le 1 - \frac{d^2 - 1}{k(d-1) + d^2 - 1} \sim 1 - \frac{1}{k}$ 



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- ▶ Optimal parallel ⇒ delayed input-state



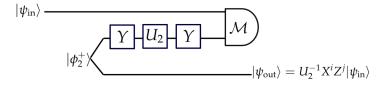
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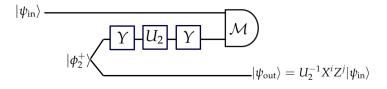
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# Qubit adaptive circuit

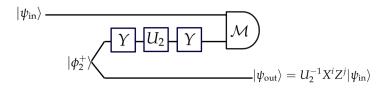


### Qubit adaptive circuit



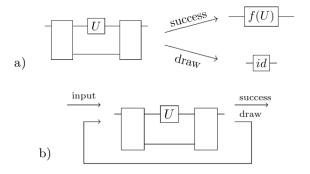
If we fail, we "destroy" the unknown input state... then we cannot re-iterate this protocol...

### Qubit adaptive circuit

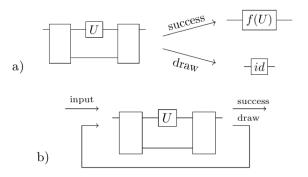


But well... if we can use  $U_2$  once more: apply  $X^{-i}Z^{-j}U_2$  to recover the input state!

### Success or draw



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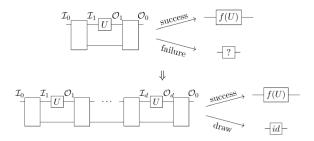
With k uses, this approach leads to a success probability of

$$p_s = 1 - (1 - p_{\mathsf{draw}})^k$$

#### Success or draw

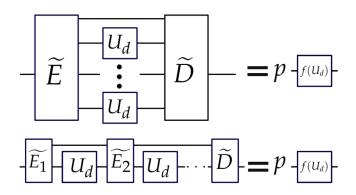
#### **Theorem**

Success or draw is always possible!

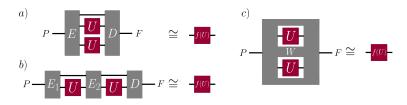


Q. Dong, MTQ, A. Soeda, M. Murao PRL (2021)

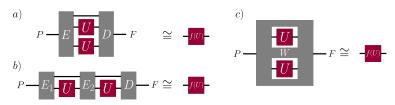
# Arbitrary functions $f(U_d)$



▶ Deterministic non-exact is also interesting!

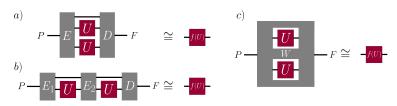


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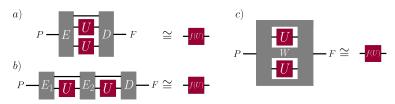
Average channel fidelity (over the Haar measure)

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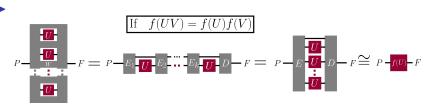
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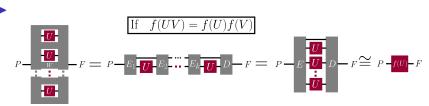


- Average channel fidelity (over the Haar measure)
- Maximal white noise visibility
- ► Parallel inversion = Parallel transposition = Unitary estimation

  MTQ, D. Ebler, Quantum 2022



A. Bisio, G.M. D'Ariano, P. Perinotti, M. Sedlak Physics Letters A (2014)

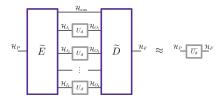


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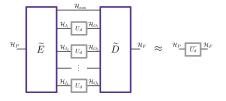
• e.g., 
$$(UV)^* = U^*V^*$$
, and cloning

▶ When  $f: \mathcal{SU}(d) \to \mathcal{SU}(d)$  and f(UV) = f(U)f(V), we know "everything".

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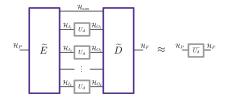


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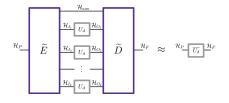
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- Step 3: use SDP duality and group theoretic results on Young-Tableau to prove  $F \leq \frac{k+1}{d(d-1)}$

arXiv:2206 00107

D. Ebler, M. Horodecki, M. Marciniak, T. Młynik, MTQ, M. Studziński



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  - D. Ebler, M. Horodecki, M. Marciniak, T. Młynik, MTQ, M. Studziński
- ▶ Step 4: prove that if k < d 1, p = 0

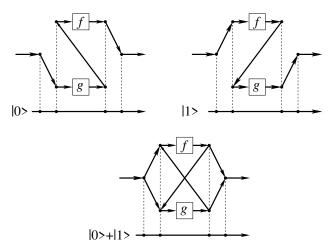


More general superchannels?

Can we go beyond sequential quantum circuits?

# More general superchannels?

### Quantum Switch:



Quantum computations without definite causal structure G. Chiribella, G. M. D'Ariano, P. Perinotti, B. Valiron PRA 2013



# More general superchannels?

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## More general superchannels?

$$\widetilde{\widetilde{S}}(\widetilde{\Lambda_1}\otimes\widetilde{\Lambda_2})=\widetilde{\Lambda_{out}}$$

Process Matrices! (May have an indefinite causal order)

- G. Chiribella, G. M. D'Ariano, P. Perinotti, and B. Valiron, PRA (2013)
- O. Oreshkov, F. Costa, and Č. Brukner, Nature Communications (2012)

 Process matrices with indefinite causal order are not (explicitly) forbidden by quantum mechanics

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- Sometimes useless
- Sometimes useful (but with limited power)

# Summary of results

Deterministic unitary inversion:  $U^{\otimes k} \stackrel{\approx}{\mapsto} U^{-1}$ 

	Parallel	Sequential	General
k=1	$F = \frac{2}{d^2}$	$F = \frac{2}{d^2}$	$F = \frac{2}{d^2}$
d=2	$F = 1 - \sin^2\left(\frac{\pi}{k+3}\right)$	?	?
(d,k)	$  M_{\mathrm{est}}  $	?	?
$k \to \infty$	$F \approx 1 - \frac{1}{k^2}$	$1 - \frac{1}{e^k} \le F \le ?$	?
Estimation	$F_{\rm par} = F_{\rm est}$	N/A	N/A
PBT	?	N/A	N/A
$k \le d-1$	$F = \frac{k+1}{d^2}$	$F = \frac{k+1}{d^2}$	$F = \frac{k+1}{d^2}$
d = 2, k = 4	$F = 1 - \sin^2\left(\frac{\pi}{7}\right)$	F=1	F=1
d = 2, k = 2	$F = 1 - \sin^2\left(\frac{\pi}{5}\right)$	$F = \frac{3}{4}$	$F \approx 0.8249$
d = 3, k = 2	$F = \frac{1}{3}$	$F = \frac{1}{3}$	$F = \frac{1}{3}$

# Summary of results

Deterministic unitary inversion: $U^{\otimes k} \stackrel{\approx}{\mapsto} U^{-1}$			Dete	erministic unitary tran	sposition: $U^{\otimes k} \stackrel{\approx}{\mapsto}$	$\cdot U^T$	
	Parallel	Sequential	General		Parallel	Sequential	General
k = 1	$F = \frac{2}{d^2}$	$F = \frac{2}{d^2}$	$F = \frac{2}{d^2}$	k = 1	$F = \frac{2}{d^2}$	$F = \frac{2}{d^2}$	$F = \frac{2}{d^2}$
d = 2	$F = 1 - \sin^2\left(\frac{\pi}{k+3}\right)$	?	?	d = 2	$F = 1 - \sin^2\left(\frac{\pi}{k+3}\right)$	?	?
(d,k)	$  M_{est}  $	?	?	(d,k)	$  M_{est}  $	?	?
$k \to \infty$	$F \approx 1 - \frac{1}{k^2}$	$1 - \frac{1}{e^k} \le F \le ?$	?	$k \to \infty$	$F \approx 1 - \frac{1}{k^2}$	$1 - \frac{1}{e^k} \le F \le ?$	?
Estimation	$F_{\text{par}} = F_{\text{est}}$	N/A	N/A	Estimation	$F_{\text{par}} = F_{\text{est}}$	N/A	N/A
PBT	?	N/A	N/A	PBT	?	N/A	N/A
$k \le d-1$	$F = \frac{k+1}{d^2}$	$F = \frac{k+1}{d^2}$	$F = \frac{k+1}{d^2}$	$k \le d-1$	$F = \frac{k+1}{d^2}$	?	?
d = 2, k = 4	$F = 1 - \sin^2(\frac{\pi}{7})$	F = 1	F = 1	d = 2, k = 4	$F = 1 - \sin^2\left(\frac{\pi}{7}\right)$	F = 1	F = 1
d = 2, k = 2	$F = 1 - \sin^2\left(\frac{\pi}{5}\right)$	$F = \frac{3}{4}$	$F \approx 0.8249$	d = 2, k = 2	$F = 1 - \sin^2\left(\frac{\pi}{5}\right)$	$F = \frac{3}{4}$	$F \approx 0.8249$
d = 3, k = 2	$F = \frac{1}{3}$	$F = \frac{1}{3}$	$F = \frac{1}{3}$	d = 3, k = 2	$F = \frac{1}{3}$	$F \approx 0,4074$	F = 0.4349

# Summary of results

Deterministic	unitary	inversion:	$U^{\otimes k}$	$\stackrel{\sim}{\mapsto}$	$U^{-}$	1

	Parallel	Sequential	General			
k = 1	$F = \frac{2}{d^2}$	$F = \frac{2}{d^2}$	$F = \frac{2}{d^2}$			
d = 2	$F = 1 - \sin^2\left(\frac{\pi}{k+3}\right)$	?	?			
(d, k)	$  M_{est}  $	?	?			
$k \to \infty$	$F \approx 1 - \frac{1}{k^2}$	$1 - \frac{1}{e^k} \le F \le ?$	?			
Estimation	$F_{\text{par}} = F_{\text{est}}$	N/A	N/A			
PBT	?	N/A	N/A			
$k \le d-1$	$F = \frac{k+1}{d^2}$	$F = \frac{k+1}{d^2}$	$F = \frac{k+1}{d^2}$			
d = 2, k = 4	$F = 1 - \sin^2\left(\frac{\pi}{7}\right)$	F = 1	F = 1			
d = 2, k = 2	$F = 1 - \sin^2\left(\frac{\pi}{5}\right)$	$F = \frac{3}{4}$	$F \approx 0.8249$			
d = 3, k = 2	$F = \frac{1}{3}$	$F = \frac{1}{3}$	$F = \frac{1}{3}$			

Probabilistic unitary inversion:  $U^{\otimes k} \mapsto p \, U^{-1}$ 

	Parallel	Sequential	General
k = 1	$p = \frac{1}{d^2}$	$p = \frac{1}{d^2}$	$p = \frac{1}{d^2}$
d = 2	$p = 1 - \frac{3}{k+3}$	?	?
(d, k)	?	?	?
$k \to \infty$	$p \approx 1 - \frac{1}{k}$	$1 - \frac{1}{e^k} \le p \le ?$	?
Store-retrieve	?	N/A	N/A
PBT	?	N/A	N/A
$k \le d-1$	?	?	?
d = 2, k = 4	$p = \frac{4}{7}$	p = 1	p = 1
d = 2, k = 2	$p = \frac{2}{5}$	$p = \frac{3}{7}$	$p \approx 4/9$
d = 3, k = 2	$p \approx \frac{1}{9}$	$p \approx \frac{1}{9}$	$p \approx \frac{1}{9}$

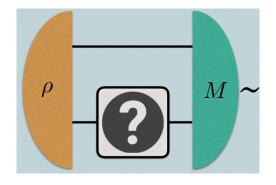
Deterministic unitary transposition: $U^{\otimes \kappa} \mapsto U^{T}$				
	Parallel	Sequential	General	
k = 1	$F = \frac{2}{d^2}$	$F = \frac{2}{d^2}$	$F = \frac{2}{d^2}$	
d = 2	$F = 1 - \sin^2\left(\frac{\pi}{k+3}\right)$	?	?	
(d,k)	$  M_{est}  $	?	?	
$k \to \infty$	$F \approx 1 - \frac{1}{k^2}$	$1 - \frac{1}{e^k} \le F \le ?$	?	
Estimation	$F_{\text{par}} = F_{\text{est}}$	N/A	N/A	
PBT	?	N/A	N/A	
$k \le d-1$	$F = \frac{k+1}{d^2}$	?	?	
d = 2, k = 4	$F = 1 - \sin^2\left(\frac{\pi}{7}\right)$	F = 1	F = 1	
d = 2, k = 2	$F = 1 - \sin^2(\frac{\pi}{5})$	$F = \frac{3}{4}$	$F \approx 0.8249$	
d = 3, k = 2	$F = \frac{1}{3}$	$F \approx 0,4074$	F = 0.4349	

Probabilistic unitary transposition: $U^{\otimes \kappa} \mapsto p U^{T}$					
	Parallel	Sequential	General		
k = 1	$p = \frac{1}{d^2}$	$p = \frac{1}{d^2}$	$p = \frac{1}{d^2}$		
d = 2	$p = 1 - \frac{3}{k+3}$	?	?		
(d,k)	$p = 1 - \frac{d^2-1}{k+d^2-1}$	?	?		
$k \to \infty$	$p \approx 1 - \frac{1}{k}$	$1 - \frac{1}{e^k} \le p \le ?$	?		
Store-retrieve	$p_{\text{par}} = p_{\text{s-r}}$	N/A	N/A		
PBT	$p_{\rm par} = p_{\rm pbt}$	N/A	N/A		
$k \le d - 1$	$p = 1 - \frac{d^2 - 1}{k + d^2 - 1}$	?	?		
d = 2, k = 4	$p = \frac{4}{7}$	p = 1	p = 1		
d = 2, k = 2	$p = \frac{2}{5}$	$p = \frac{3}{7}$	$p \approx 4/9$		
d = 3, k = 2	$p = \frac{2}{10}$	$ppproxrac{2}{9}$	$p \approx \frac{2}{8}$		

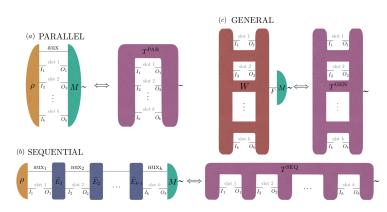
# Measuring quantum operations

How can one perform a measurement on a quantum operation?

# Measuring quantum operations



## Measuring quantum operations



- J. Bavaresco, M. Murao, MTQ PRL (2021)
- J. Bavaresco, M. Murao, MTQ J. Math. Phys. (2022)

► Two unitaries ⇒ parallel is optimal G. Chiribella, G. M. D'Ariano, P. Perinotti PRL (2008)

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- lackbox Uniform unitaries + a group structure  $\Longrightarrow$  parallel is optimal
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- ▶ Non-symmetric scenarios? Non-unitary quantum channels?

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  G. Chiribella, G. M. D'Ariano, P. Perinotti PRL (2008)
- lacktriangle Uniform unitaries + a group structure  $\implies$  parallel is optimal
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  - J. Bavaresco, M. Murao, MTQ, J. Math. Phys. (2022)
- ▶ Non-symmetric scenarios? Non-unitary quantum channels?
- ► Almost always, we have a strict hierarchy!
  - J. Bavaresco, M. Murao, MTQ PRL (2021)
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- Power and limitation of indefinite causality
- ► Can we have deterministic exact unitary inversion with finite sequential circuits?
- ▶ Practical and realistic implementations

# Thank you!





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Qingxiuxiong Dong



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Akihito Soeda



Jessica Bavaresco



Michał Studziński



Michał Horodecki



Tomasz Młynik



Daniel Ebler

# Thank you!

