Parallel, sequential, and non-causal strategies for transforming unitary operations and discriminating quantum channel via a higher-order approach

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## Quantum state transformations

$$
\left|\psi_{\mathrm{in}}\right\rangle \mapsto\left|\psi_{\text {out }}\right\rangle
$$

Quantum state transformations

$$
\begin{gathered}
\left|\psi_{\text {in }}\right\rangle \mapsto\left|\psi_{\text {out }}\right\rangle=U\left|\psi_{\text {in }}\right\rangle \\
U^{\dagger} U=I
\end{gathered}
$$

## Quantum state transformations

$\rho_{\text {in }} \longmapsto \rho_{\text {out }}$

Quantum state transformations

$$
\begin{gathered}
\rho_{\text {in }} \mapsto \rho_{\text {out }}=\widetilde{\Lambda}\left(\rho_{\text {in }}\right) \\
\widetilde{\Lambda} \text { is CPTP }
\end{gathered}
$$

## Quantum operation transformations

Can we transform quantum operations??

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U_{\text {in }} \mapsto U_{\text {out }}
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Quantum operation transformations

Can we transform quantum operations??
$U_{\text {in }} \mapsto U_{\text {out }}$
$\widetilde{\Lambda_{\text {in }}} \mapsto \widetilde{\Lambda_{\text {out }}}$

## "Quantum" unitary inversion

$$
U_{d} \mapsto U_{d}^{-1}
$$

The universal/unknown paradigm

$$
\sigma_{Z}=\left[\begin{array}{cc}
1 & 0 \\
0 & -1
\end{array}\right]=\sigma_{Z}^{-1}
$$

What do we want?

## What do we want?

Ideally...
Something like this:


Phys. Rev. Research (2019)
J. Miyazaki, A. Soeda, and M. Murao

## What do we want?

- Universal (also works for "unknown" d-dimensional unitary)


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## What do we want?

- Universal (also works for "unknown" d-dimensional unitary)
- Exact
- Possible?
- Optimal average fidelity: $F_{\max }=\frac{2}{d^{2}}$
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- $F_{\text {max }}<1 \Longrightarrow$ Impossible...

What do we want?

Probabilistic heralded?

What do we want?

## Probabilistic heralded? For qubits, Possible!

## Explicit construction



## Explicit construction



## Delayed input state protocols



## We want more!

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- Higher-order operations and supermaps!


## Superchannels



The most general quantum superchannel?

## Superchannels



$$
\widetilde{\widetilde{S}}:\left[\mathcal{L}\left(\mathcal{H}_{2}\right) \rightarrow \mathcal{L}\left(\mathcal{H}_{3}\right)\right] \rightarrow\left[\mathcal{L}\left(\mathcal{H}_{1}\right) \rightarrow \mathcal{L}\left(\mathcal{H}_{4}\right)\right]
$$

- $\widetilde{\widetilde{S}}$ is a linear supermap


## Superchannels



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- $\widetilde{\widetilde{S}}$ maps valid channels into valid channels (TP preserving, CP preserving)


## Superchannels



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- $\widetilde{\widetilde{S}}$ is a linear supermap
- $\widetilde{\widetilde{S}}$ maps valid channels into valid channels (TP preserving, CP preserving)
- $\widetilde{\widetilde{S}}$ may be applied into part of channel (completely CP preserving)


## Superchannels



$$
\widetilde{\widetilde{S}}\left(\widetilde{\Lambda_{\text {in }}}\right)=\operatorname{tr}_{A}\left(\widetilde{D} \circ\left(\widetilde{\Lambda_{\text {in }}} \otimes \widetilde{I_{A}}\right) \circ \widetilde{E}\right)
$$

G. Chiribella, G. M. D'Ariano, and P. Perinotti EPL (2008)
K. Życzkowski J. Phys. A 41, 355302-23 (2008)
G. Gutoski and J. Watrous Proceedings of STOC (2007)

## Superchannels



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- How to represent such mathematical objects?


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- Maps $\widetilde{\Lambda}: \mathcal{L}\left(\mathcal{H}_{2}\right) \rightarrow \mathcal{L}\left(\mathcal{H}_{3}\right)$ become matrices $\Lambda \in \mathcal{L}\left(\mathcal{H}_{2} \otimes \mathcal{H}_{3}\right)$.


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- Maps $\widetilde{S}$ become matrices $S \in \mathcal{L}\left(\mathcal{H}_{1} \otimes \mathcal{H}_{2} \otimes \mathcal{H}_{3} \otimes H_{4}\right)$


## Superchannels


$S \in \mathcal{L}\left(\mathcal{H}_{1} \otimes \mathcal{H}_{2} \otimes \mathcal{H}_{3} \otimes H_{4}\right)$ is a superchannel iff

$$
\begin{aligned}
S & \geq 0 \\
\operatorname{tr}_{4}(S) & =\operatorname{tr}_{34}(S) \otimes \frac{l_{4}}{d_{4}} \\
\operatorname{tr}_{234}(S) & =\operatorname{tr}_{1234}(S) \otimes \frac{l_{2}}{d_{2}} \\
\operatorname{tr}(S) & =d_{1} d_{3}
\end{aligned}
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Affine and positive semidefinite constraints $\Longrightarrow$ SDP!!

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(Quantum combs, channel with memory, quantum strategy, quantum superchannels with multiple inputs)

Results

$\widetilde{E_{1}}-\sqrt[U_{d}]{ }-\widetilde{E_{2}}-U_{d} \cdot \cdots, \widetilde{D}=p \sqrt{U_{d}^{-1}}$

- Parallel $(d=2): \quad p=1-\frac{3}{k+3}$


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$1-\frac{1}{k} \sim 1-\frac{d^{2}-1}{\left\lfloor\frac{k}{d-1}\right\rfloor+d^{2}-1} \leq p \leq 1-\frac{d^{2}-1}{k(d-1)+d^{2}-1} \sim 1-\frac{1}{k}$


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- Optimal parallel $\Longrightarrow$ delayed input-state
- Sequential $(k<d-1): p=0$
- Sequential $(k \geq d-1): p \geq 1-\left(1-\frac{1}{d^{2}}\right)^{\left\lceil\frac{k+2-d}{d}\right\rceil} \sim 1-\frac{1}{e^{k}}$ MTQ, Q. Dong, A. Shimbo, A. Soeda, M. Murao PRA (2019), PRL (2019)


## Qubit adaptive circuit



## Qubit adaptive circuit



If we fail, we "destroy" the unknown input state... then we cannot re-iterate this protocol...

## Qubit adaptive circuit



But well. . .
if we can use $U_{2}$ once more:
apply $X^{-i} Z^{-j} U_{2}$ to recover the input state!

## Success or draw

a)

b)


## Success or draw



With $k$ uses, this approach leads to a success probability of

$$
p_{s}=1-\left(1-p_{\text {draw }}\right)^{k}
$$

## Success or draw

Theorem
Success or draw is always possible!

Q. Dong, MTQ, A. Soeda, M. Murao

PRL (2021)

## Arbitrary functions $f\left(U_{d}\right)$


$\sqrt{\widetilde{E_{1}}} \sqrt{u_{d}} \sqrt{E_{2}} \sqrt{U_{d}} \cdots \widetilde{D}=p-\left\{\left(u_{d}\right)\right.$

## Deterministic protocols?

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## Deterministic protocols?

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- Average channel fidelity (over the Haar measure)
- Maximal white noise visibility
- Parallel inversion = Parallel transposition $=$ Unitary estimation MTQ, D. Ebler, Quantum 2022


## Solving the homomorphic case


A. Bisio, G.M. D'Ariano, P. Perinotti, M. Sedlak Physics Letters A (2014)

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A. Bisio, G.M. D'Ariano, P. Perinotti, M. Sedlak Physics Letters A (2014)

- e.g., $(U V)^{*}=U^{*} V^{*}$, and cloning


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arXiv:2206.00107
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- Step 4: prove that if $k<d-1, p=0$

More general superchannels?

Can we go beyond sequential quantum circuits?

## More general superchannels?

Quantum Switch:


Quantum computations without definite causal structure
G. Chiribella, G. M. D'Ariano, P. Perinotti, B. Valiron

PRA 2013

More general superchannels?

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Process Matrices! (May have an indefinite causal order)
G. Chiribella, G. M. D'Ariano, P. Perinotti, and B. Valiron, PRA (2013)
O. Oreshkov, F. Costa, and Č. Brukner, Nature Communications (2012)

## Process matrices

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- Upper bounds and ansatz
- Sometimes useless
- Sometimes useful (but with limited power)


## Summary of results

Deterministic unitary inversion: $U^{\otimes k} \stackrel{\approx}{\leftrightarrows} U^{-1}$

|  | Parallel | Sequential | General |
| :---: | :---: | :---: | :---: |
| $k=1$ | $F=\frac{2}{d^{2}}$ | $F=\frac{2}{d^{2}}$ | $F=\frac{2}{d^{2}}$ |
| $d=2$ | $F=1-\sin ^{2}\left(\frac{\pi}{k+3}\right)$ | $?$ | $?$ |
| $(d, k)$ | $\left\\|M_{\text {est }}\right\\|$ | $?$ | $?$ |
| $k \rightarrow \infty$ | $F \approx 1-\frac{1}{k^{2}}$ | $1-\frac{1}{e^{k}} \leq F \leq ?$ | $?$ |
| Estimation | $F_{\text {par }}=F_{\text {est }}$ | N/A | N/A |
| PBT | $?$ | N/A | N/A |
| $k \leq d-1$ | $F=\frac{k+1}{d^{2}}$ | $F=\frac{k+1}{d^{2}}$ | $F=\frac{k+1}{d^{2}}$ |
| $d=2, k=4$ | $F=1-\sin ^{2}\left(\frac{\pi}{7}\right)$ | $F=1$ | $F=1$ |
| $d=2, k=2$ | $F=1-\sin ^{2}\left(\frac{\pi}{5}\right)$ | $F=\frac{3}{4}$ | $F \approx 0.8249$ |
| $d=3, k=2$ | $F=\frac{1}{3}$ | $F=\frac{1}{3}$ | $F=\frac{1}{3}$ |

## Summary of results

Deterministic unitary inversion: $U^{\otimes k} \stackrel{\approx}{\leftrightarrows} U^{-1}$
Deterministic unitary transposition: $U^{\otimes k} \stackrel{\approx}{\rightleftarrows} U^{T}$

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| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $k=1$ | $F=\frac{2}{d^{2}}$ | $F=\frac{2}{d^{2}}$ | $F=\frac{2}{d^{2}}$ | $k=1$ | $F=\frac{2}{d^{2}}$ | $F=\frac{2}{d^{2}}$ | $F=\frac{2}{d^{2}}$ |
| $d=2$ | $F=1-\sin ^{2}\left(\frac{\pi}{k+3}\right)$ | ? | ? | $d=2$ | $F=1-\sin ^{2}\left(\frac{\pi}{k+3}\right)$ | ? | ? |
| (d, $k$ ) | $\left\\|M_{\text {est }}\right\\|$ | ? | ? | $(d, k)$ | $\left\|M_{\text {est }}\right\| \mid$ | ? | ? |
| $k \rightarrow \infty$ | $F \approx 1-\frac{1}{k^{2}}$ | $1-\frac{1}{e^{k}} \leq F \leq$ ? | ? | $k \rightarrow \infty$ | $F \approx 1-\frac{1}{k^{2}}$ | $1-\frac{1}{e^{k}} \leq F \leq$ ? | ? |
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| PBT | ? | N/A | N/A | PBT | ? | N/A | N/A |
| $k \leq d-1$ | $F=\frac{k+1}{d^{2}}$ | $F=\frac{k+1}{d^{2}}$ | $F=\frac{k+1}{d^{2}}$ | $k \leq d-1$ | $F=\frac{k+1}{d^{2}}$ | ? | ? |
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Deterministic unitary transposition: $U^{\otimes k} \stackrel{\approx}{\rightleftarrows} U^{T}$

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Probabilistic unitary transposition: $U^{\otimes k} \mapsto p U^{T}$

|  | Parallel | Sequential | General |
| :---: | :---: | :---: | :---: |
| $k=1$ | $p=\frac{1}{d^{2}}$ | $p=\frac{1}{d^{2}}$ | $p=\frac{1}{d^{2}}$ |
| $d=2$ | $p=1-\frac{3}{k+3}$ | $?$ | $?$ |
| $(d, k)$ | $p=1-\frac{d^{2}-1}{k+d^{2}-1}$ | $?$ | $?$ |
| $k \rightarrow \infty$ | $p \approx 1-\frac{1}{k}$ | $1-\frac{1}{e^{k}} \leq p \leq ?$ | $?$ |
| Store-retrieve | $p_{\text {par }}=p_{\mathrm{s}-\mathrm{r}}$ | $\mathrm{N} / \mathrm{A}$ | $\mathrm{N} / \mathrm{A}$ |
| PBT | $p_{\text {par }}=p_{\mathrm{pbt}}$ | $\mathrm{N} / \mathrm{A}$ | $\mathrm{N} / \mathrm{A}$ |
| $k \leq d-1$ | $p=1-\frac{d^{2}-1}{k+d^{2}-1}$ | $?$ | $?$ |
| $d=2, k=4$ | $p=\frac{4}{7}$ | $p=1$ | $p=1$ |
| $d=2, k=2$ | $p=\frac{2}{5}$ | $p=\frac{3}{7}$ | $p \approx 4 / 9$ |
| $d=3, k=2$ | $p=\frac{2}{10}$ | $p \approx \frac{2}{9}$ | $p \approx \frac{2}{8}$ |

## Measuring quantum operations

## How can one perform a measurement on a quantum operation?

Measuring quantum operations


## Measuring quantum operations

(c) GENERAL
(a) PARALLEL

(b) SEQUENTIAL

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## Quantum channel discrimination

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- Non-symmetric scenarios? Non-unitary quantum channels?
- Almost always, we have a strict hierarchy!
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- Practical and realistic implementations


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